

SPARROW+CP3 and SPARROWTORISS

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Abstract—SPARROW+CP3 and SPARROWTORISS are using as a first step the preprocessor CP3 to simplify the formula in a way that is beneficial for SLS solvers. SPARROW+CP3 then uses the solver SPARROW to solve the simplified problem. SPARROWTORISS is first trying to solve the problem with Sparrow, limiting its execution to $5 \cdot 10^8$ flips and then passes the assignment found to the CDCL solver RISS3G, which uses this information for initialization and then tries to solve the problem. The solver RISS3G combines the improved Minisat-style solving engine of GLUCOSE 2.2 with a state-of-the-art preprocessor COPROCESSOR and adds further modifications to the search process. The SLS solver SPARROW is an improved version of SPARROW 2011

I. INTRODUCTION

SLS solvers showed remarkable performance on the satisfiable crafted problems in the competitions from the last years. Motivated by this results we have analyzed in [1] the utility of different preprocessing techniques for the SLS solver SPARROW. The best found technique together with SPARROW represents the basis of our solver SPARROW+CP3.

As SPARROW is not able to prove the unsatisfiability of a problem we have decided to append a CDCL solver to SPARROW+CP3, namely RISS3G after limiting the execution of SPARROW to $5 \cdot 10^8$ flips. The CDCL solver RISS3G uses the MINISAT search engine [2], more specifically the extensions added in GLUCOSE 2.2 [3], [4]. Furthermore, RISS3G is equipped with the preprocessor COPROCESSOR(CP3) [5], that implements most of the recently published formula simplification techniques, ready to be used as inprocessing as well.

II. MAIN TECHNIQUES

SPARROW is a clause weighting SLS solvers that uses promising variables and probability distribution based selection heuristics. It is described in detail in [6]. Compared to the original version, the one submitted here is updating weights of unsatisfied clauses in every step where no promising variable can be found.

The built-in preprocessor CP3 has been ported from CO-PROCESSOR 2 and supports the following simplification techniques: Unit Propagation, Subsumption, Strengthening (also called self-subsuming resolution) – where for small clauses all subsuming resolvents can be produced, (Bounded) Variable Elimination (BVE) [7] combined with Blocked Clause Elimination (BCE) [8], (Bounded) Variable Addition (BVA) [9], Probing [10], Covered Clause Elimination [11], Hidden Tautology Elimination [12], Equivalent Literal Elimination [13], Unhiding (Unhide) [14], Add Binary Resolvents [15], at-most-one rewriting [16], [17], a 2SAT algorithm [18], and a walksat

implementation [19]. The preprocessor furthermore supports parallel subsumption, strengthening and variable elimination, which is described in [20].

RISS3G uses GLUCOSE 2.2 as main search engine – the version used in SPARROWTORISS just replaces the internal preprocessor with CP3.

The combination of the SPARROW and RISS3G, called SPARROWTORISS, does not simply execute the two solvers after each other, but also forwards information from the SLS solver to the CDCL solver: when SPARROW terminates, it outputs its last full assignment in chronological order (i.e. the oldest variable first), which is used to initialize the phase saving of RISS3G, such that the first decisions of RISS3G follow this assignment. In a brief empirical evaluation this communication turned out to be useful. The solvers are also able to forward the information about the age of the variables in the SLS search. This data could be used to initialize the activities of the variables inside RISS3G. However, this feature is not enabled in the used configuration.

III. MAIN PARAMETERS

SPARROW is using the same parameters as SPARROW 2011.

The configuration of CP3 has been tuned for SPARROW in [1] on the SAT Challenge 2012 satisfiable hard combinatorial benchmarks.

The main parameters of RISS3G control how the formula simplification of CP3 is executed. The configuration of CP3 has been tuned for GLUCOSE 2.2 in [1] on the SAT Challenge 2012 application benchmark. The final setup of the preprocessor inside RISS3G uses the following techniques: UP, SUB+STR (producing all resolvents for ternary clauses), Unhide without *hidden literal elimination* [14] and 5 iterations, BVE without on the fly BCE and BVA with a small number of 120000 steps.

For SPARROWTORISS it can be chosen whether to forward the last assignment, or the activity information.

IV. IMPLEMENTATION DETAILS

SPARROW is implemented in C. The solver RISS3G is build on top of MINISAT 2.2 and GLUCOSE 2.2. Furthermore, we integrated COPROCESSOR into the system, allowing inprocessing techniques to be executed during search – however, this feature is not used in the competition. All solvers have been compiled with the GCC C++ compiler as 64 bit binaries.

V. AVAILABILITY

The source code of RISS3G (including CP3) is available at tools.computational-logic.org for research purposes.

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